Universality of one-dimensional reversible and number-conserving cellular automata

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Outline of the Talk

- Objective: Finding universal 1-d reversible and number-conserving CAs (RNCCAs).
- We give a method of converting a 2-neighbor s-state reversible partitioned CA (RPCA) into a 4-neighbor 4s-state RNCCA.
- Since there is a universal 2-neighbor 24-state RPCA [Morita, 2011], we can obtain a computa-tionally universal 4-neighbor 96-state RNCCA.

Contents

1. Preliminaries

- Number-Conserving CAs (NCCAs)
- Reversible NCCAs (RNCCAs)
- Partitioned CAs (PCAs)
- 2. Converting a reversible PCA (RPCA) into an RNCCA

1. Preliminaries

- Number-Conserving CAs (NCCAs)
- Reversible NCCAs (RNCCAs)
- Partitioned CAs (PCAs)

Number-Conserving CA (NCCA)

- An abstract spatiotemporal model having a property similar to the conservation laws in physics.
 - Each cell has an integer value.
 - Their sum in a configuration is conserved.

Sum

$$t \quad \cdots \quad 0 \quad 0 \quad 5 \quad 2 \quad 11 \quad 3 \quad 7 \quad 0 \quad 0 \quad \cdots \quad 28$$

$$t+1$$
 \cdots 0 0 3 4 9 4 5 3 0 \cdots 28

 So far, several definitions and characterization of NCCAs have been given [Boccara, Fuks, 2000], [Durand, et al., 2003].

Several Notions on NCCAs

- Periodic-number-conserving
- Finite-number-conserving
- Number-conserving (for infinite configurations)

Durand, Formenti, Róka (2003) proved the above are all equivalent.

• Note: In our proof, we use the definition of *finitenumber-conserving* CAs.

Finite-Number-Conserving CAs

Definition [Durand, et al., 2003]

- Let A be a 1-d CA whose state set is $Q = \{0, \dots, s-1\}.$
- Let F be the global function of A.
- Let Conf_{fin}(Q) be the set of all finite configurations over Q.

A is called *finite-number-conserving*, if the following condition holds.

$$\forall \alpha \in \mathsf{Conf}_{\mathsf{fin}}(Q) : \sum_{x \in \mathbb{Z}} \alpha(x) = \sum_{x \in \mathbb{Z}} F(\alpha)(x)$$

Reversible Cellular Automaton (RCA)

- It is a CA whose global function is one-to-one.
- Hence, there is no pair of configurations that go to the same configuration.



• In the following, we investigate *reversible NCCAs* (RNCCAs).

Past Studies on 1-D RNCCAs

- 2-neighbor (radius 1/2) case: Every RNCCA is a *shift-identity product CA*.
 [García-Ramos, 2012]
- 3-neighbor (radius 1) case:
 Some RNCCAs show nontrivial behaviors.
 [Imai, Martin, Saito, 2012]

2-Neighbor (Radius 1/2) RNCCAs

They are all shift-identity product CAs (SIPCAs).
 [García-Ramos, 2012]



In a SIPCA, signals do not interact each other.Hence, it cannot be computationally universal.

Past Studies on 1-D RNCCAs

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3-Neighbor (Radius 1) RNCCAs

• Some RNCCAs show nontrivial behaviors.

[Imai, Martin, Saito, 2012]



• In these CAs, signals can interact each other.

• It is not known whether there is a universal one.

Studies on 1-D RNCCAs

- 2-neighbor (radius 1/2) case: Every RNCCA is a *shift-identity product CA*.
 [García-Ramos, 2012]
- 3-neighbor (radius 1) case:
 Some RNCCAs show nontrivial behaviors.
 [Imai, Martin, Saito, 2012]
- 4-neighbor (radius 3/2) case: In this talk, we will show there is a *universal* RNCCA.

Partitioned Cellular Automaton (PCA)

1-d 3-neighbor PCA



f: local function

• To construct an RCA, it is sufficient to give a PCA whose local function f is one-to-one.

1-D 2-Neighbor PCA

 It is a special case of a 3-neighbor PCA where the left state set is a singleton.



A 2-Neighbor RPCA Can Simulate a 3-Neighbor RPCA [Morita, 1992]



3-Neighbor PCA P₃



2-Neighbor PCA P₂

Computationally Universal 1-D Reversible CAs

- On infinite configurations:
 24-state 2-neighbor RPCA
- On finite configurations:
 98-state 3-neighbor RPCA

[Morita, 2011]

[Morita, 2007]

A 24-State 1-D 2-Neighbor RPCA That Simulates Any Cyclic Tag System

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A 24-State 1-D 2-Neighbor RPCA That Simulates Any Cyclic Tag System (Colored version)



How can we convert it into an RNCCA?

2. Converting a reversible PCA (RPCA) into an RNCCA

Main Result

Lemma

• Let $P = (\mathbb{Z}, (C, R), (0, -1), f)$ be a 1-d RPCA.

• We can construct a 1-d RNCCA

 $A = (\mathbb{Z}, \tilde{Q}, (-2, -1, 0, 1), \tilde{f}, 0)$ that simulates P, such that $\tilde{Q} = \{0, 1, \dots, 4|C| \cdot |R| - 1\}.$

Proposition [Morita, 2011] There is a computationally universal 1-d 2-neighbor 24-state RPCA.

Theorem There is a computationally universal 1-d 4-neighbor 96-state RNCCA.

How to Construct an RNCCA from an RPCA — A Proof Outline of the Lemma —

Given: A 1-d 2-neighbor s-state RPCA $P = (\mathbb{Z}, (C, R), (0, -1), f)$

Construct: A 1-d 4-neighbor 4*s*-state RNCCA $A = (\mathbb{Z}, \tilde{Q}, (-2, -1, 0, 1), \tilde{f}, 0)$

Example:

If $C = \{Y, N, P, M\}$ and $R = \{y, n, +, -, *, /\}$, then $\tilde{Q} = \{0, ..., 95\}$, and \tilde{f} is determined as below.

Heavy and Light Particles for the States in C and R

 $P = (\mathbb{Z}, (C, R), (0, -1), f)$

• If $C = \{Y, N, P, M\}$ and $R = \{y, n, +, -, *, /\}$, then prepare the following 4 sets of integers: $\hat{C} = \{0, 12, 24, 36\}, \ \hat{R} = \{0, 1, 2, 3, 4, 5\}$ $\check{C} = \{84, 72, 60, 48\}, \ \check{R} = \{11, 10, 9, 8, 7, 6\}$ Elements of $\hat{C} \cup \check{C}$ are heavy particles. Elements of $\hat{R} \cup \check{R}$ are light particles.

Each number represents the *mass* of the particle.

Correspondence among $C, R, \widehat{C}, \widecheck{C}, \widehat{R},$ and \widecheck{R}															
$P = (\mathbb{Z}, (C, R), (0, -1), f)$															
C	$C = \{Y, N, P, M\}, R = \{y, n, +, -, *, /\}$														
$\hat{\varphi}_C: C \to \hat{C}, \ \check{\varphi}_C: C \to \check{C} \qquad \qquad \hat{\varphi}_R: R \to \hat{R}, \ \check{\varphi}_R: R \to \check{R}$															
С	Y	N	P	M	r	$\mid y$	n	+	—	*	/				
$\widehat{\varphi}_C(c)$	0	12	24	36	$\widehat{\varphi}_R(r)$	0	1	2	3	4	5				
$\check{\varphi}_C(c)$	84	72	60	48	$\check{arphi}_R(r)$	11	10	9	8	7	6				
$\widehat{\varphi}_C(c)$	\widehat{Y}	\widehat{N}	(\hat{P})	\widehat{M}	$\widehat{\varphi}_R(r)$	\widehat{y}	\widehat{n}	$(\widehat{+})$	$\widehat{-}$	(*)	\bigcirc				
$\check{\varphi}_C(c)$	(\check{Y})	(\check{N})	(\check{P})	(\tilde{M})	$\check{arphi}_R(r)$	(\check{y})	Ň	$(\tilde{+})$	$\tilde{-}$	(×	\bigcirc				

Hereafter, we use notations (c), (c), (r), (r).
(c), (c) and (r), (r) are complementary pairs, since (c) + (c) = 84 and (r) + (r) = 11.

Simulating RPCA P by RNCCA A (1)

 $A = (\mathbb{Z}, \tilde{Q}, (-2, -1, 0, 1), \tilde{f}, 0)$

- Each cell of A keeps one c and one r.
- A light particle r always moves rightward at the unit speed.
- A heavy particle c is stationary.



Simulating RPCA P by RNCCA A (2)

• If a complementary pair of light particles (\hat{r}, \hat{r}) meets that of heavy particles (\hat{c}, \hat{c}) , state transition of P is simulated. Then, new pairs of complementary particles are created.





• In all other cases, a light particle r simply moves rightward without interacting heavy particles.

Simulation Process of P by A $A = (\mathbb{Z}, \tilde{Q}, (-2, -1, 0, 1), \tilde{f}, 0)$



• A precise proof on number-conserving property and reversibility of A is found in the proceeding.

Concluding Remarks

• Result:

A universal 4-neighbor 96-state RNCCA exists, in spite of the strong constraints of reversibility and number-conservation.

• Open problem:

- Is there a universal 3-neighbor RNCCA?