A GRAPH THEORETIC PROBLEM ON LOCAL AREA NETWORK DESIGN

Tadashi Ae, Reiji Aibara, Wagner Chiepa Cunha and Kaori Hayashihara

Faculty of Engineering, Hiroshima University Saijo, Higashi-Hiroshima, 724 Japan

1. Introduction

In this paper we focus on a graph theoretic problem related to the design of a local area network(in short LAN) illustrated as in Fig.1, where the computing devices are separated into the user layer and the server layer. These two layers correspond to the work stations and the resources, respectively. The LAN plays a role of interconnection between two layers.

A virtual computer system for a user can be represented as a triangle of Fig.2, where the leftmost vertex means a work station. The triangle represents also a hierarchical structure of program modules, and then the LAN layer is inserted between two layers, i.e., two separated If the data between sets of program modules. two layers flows as fast as in a single computer, The realistic LAN, the problem does not occur. however, is not so fast, and then the problem to find the optimum cutset where the LAN is inserted becomes siginificant. We propose a theoretical problem arising from such background, and state an efficient algorithm to solve this.

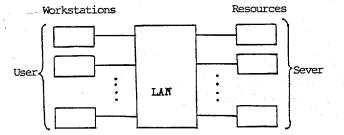


Fig.1 A conceptual LAN system

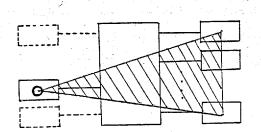


Fig.2 A virtual subsystem on LAN system

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2. Minimum Cost Cutset Problem

Since the problem is to find the minimum cost LAN which divides the system into two layers, a graph model is proposed for this purpose.

First, we introduce a kind of data flow graph that is a data flow tree where each edge has the The data flow tree value of required data flow. differs from the conventional tree graph in variety of the type of nodes. In the case of LAN the tree must include both And-type node and OR-type node, which are needed to represent the resource locking and the conditional branch in If the tree consists of programs, respectively. only AND-type nodes, the problem becomes the same as in the conventional graph[1]. An example of AND/OR tree is shown as in Fig.3 as a data flow model. The root(a double circle) is the initial point of each behavior on interconnected program modules, and the leaves (squares) are the An AND/OR tree represents a terminal modules. dynamic behavior of logical computing systems. The minimum cost problem of LAN for computing system is reduced to the problem to find the minimum cost cutset for an AND/OR tree where the data flows from the root to leaves.

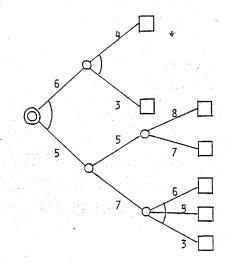


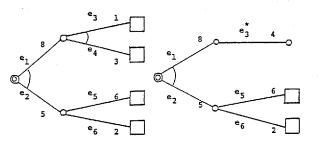
Fig.3 An example of AND/OR tree

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3. Algorithm Next, we show an algorithm to obtain the minimum cost cutset for an AND/OR tree. *COUTLINE OF ALGORITHM>* INPUT : AND/OR tree G=(V,E) with label OUTPUT : Set of Minimum Cost Cutset C ; Procedure PAND ; Input:Subtree Co whose node v has relation AND C₀ = { ei,....,ei*k}; Output: Edge e* whose label is the sum of all labels of $e_{i+p} \in C$ (p=0,...,k); Procedure POR ; Input:Subtree C_0 whose node v has relation OR $C_0 = \{ e_1, \dots, e_{i+k} \};$ Output: Edge e_i^* whose label is the maximum label of $e_{i+p} \in C_0$ (p=0,...,k); Procedure PPATH; Input: Subtree ${\rm C}_0$ whose node v has two edges C₀ = { e_i, e_j (or e^{*}_j)}: Output: Edge e^{*}₁ whose label is smaller one of two in C₀; /* main */ begin /* initialization */ C <--- E; for level := k downto 1 do while (set of nodes belong to current level) $\neq \phi do$ begin if (the objective node v has two or more edges) then begin if (node v is AND-type node) then begin PAND ; /*call procedure AND*/ $C \leftarrow C - C - C_0 + \{e^*\}$ enð else begin POR ; /*call procedure OR*/ $C \leftarrow C - C_{0} + \{e_{1}^{*}\}$ end end; PPATH ; /*call procedure PATH*/ $-C_0+\{e_i^*\}$ /* edge e; or e; is C <-- C deleted */ end end.

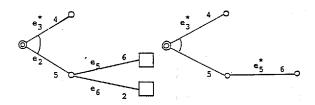
The level means the distance from the root. C_0 is the objective set of edges, and { e^{*} } is an edge reduced from C_0 by each procedure. PAND, POR or PPATH which is called by the algorithm is a lower procedure for each type node(AND, OR or UNARY), where the star-type subtree is reduced to an edge. We demonstrate the algorithm for an AND/OR tree as in Fig.4.

4. Conclusion As a result we obtain the following;

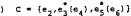


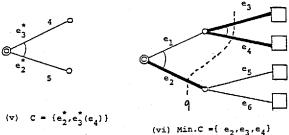
C= {e₁,e₂,e₃,e₄,e₅,e₆} (i)

(ii) C ≠ {e1,e2,e3(e4),e5,e6)



(iii) C = {e2 ,e^{*}₃(e₄),e₅,e₆} (iv)





An application of

Fig.4 ALGORITHM described in the paper

[Theorem]

There exists an O(n) algorithm to find the minimum cost cutset for an AND/OR tree. If the number of minimum cost cutsets are fixed, then the set of minimum cost cutsets is also obtained within the linear time.

An extended case of non-tree graph has been discussed. There still remain, however, some difficulties to utilize this result for realistic LAN design, but the discussion of the this paper reminds us of importance, of data flow analysis in the distributed system [2].

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